

Levitz class and punctually non-standard models of natural numbers

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of Prof. Tinko Tinchev**

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Some memories

13 years ago

Work team

This is a joint project with:

- ▶ Stefan Vatev;
- ▶ Nikolay Bazhenov (Sobolev Institute, Novosibirsk);
- ▶ Dariusz Kalociński (Polish Academy of Sciences);
- ▶ Michał Wroclawski (University of Warsaw).

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2. consider the subclass of Levitz, in which unique normal forms exist;
3. present an ad-hoc construction for punctual copies of (\mathbb{N}, S) and discuss its limitations;
4. introduce the islands and archipelago technique for overcoming these limitations.

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$$\begin{aligned} f, g \in \mathfrak{A} \Rightarrow \lambda x.[f(x) + g(x)] &\in \mathfrak{A}, \\ \lambda x.[f(x) \cdot g(x)] &\in \mathfrak{A}, \\ \lambda x.[f(x)^{g(x)}] &\in \mathfrak{A}. \end{aligned}$$

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Different terms may represent the same function, for example:

$$x \cdot (x + x) = x \cdot x + x \cdot x = (1 + 1) \cdot (x \cdot x).$$

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Note that $f \preccurlyeq g$ & $g \preccurlyeq f$ iff f and g are almost equal (they differ only on a finite set).

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For an alternative proof, one can also use Wilkie's theorem on o -minimality of $\text{Th}(\mathbb{R}^{\text{exp}})$.

Computability of identity and domination

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This question is clearly connected with the decidability of $\text{Th}(\mathbb{R}^{\text{exp}})$, which is a major open problem.

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As a simple example: the order type of $\mathbb{N}^+[x]$ is ω^ω .

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It is immaterial that we extend the functions from \mathbb{N}^+ to \mathbb{N} .

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Comparing two additive primes in \mathfrak{L} with respect to \preccurlyeq

In order to compare two additive primes $f = u_1^{f_1} u_2^{f_2} \dots u_k^{f_k}$ and $g = v_1^{g_1} v_2^{g_2} \dots v_\ell^{g_\ell}$ in multiplicative normal form:

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In order to compare $f = p_1 + p_2 + \dots + p_k$ and
 $g = q_1 + q_2 + \dots + q_\ell$ in additive normal form:

1. if f is an initial segment of g , then $f \preccurlyeq g$;
2. if g is an initial segment of f , then $g \preccurlyeq f$;
3. otherwise, compute the least i , such that $p_i \neq q_i$;
4. if $p_i \prec q_i$, then $f \prec g$;
5. if $q_i \prec p_i$, then $g \prec f$.

The existence and uniqueness of normal forms implies that the order type of $\mathfrak{L}, \preccurlyeq$ is ϵ_0 .

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f is computable if and only if $f^{\mathcal{A}}$ is computable.

But the complexity of f and $f^{\mathcal{A}}$ can be very different.

Main question

For a punctual copy \mathcal{A} we are interested in the class of primitive recursive functions, relative to \mathcal{A} :

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Our first line of work was to build specific \mathcal{A} , in which some concrete functions in $Pr(\mathcal{A})$ are not primitive recursive.

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We complete the model \mathcal{A} in the following way:

$$a_n \rightarrow \text{free}(n) \rightarrow h(\langle n, 0 \rangle)$$

Images of functions in the model \mathcal{A}

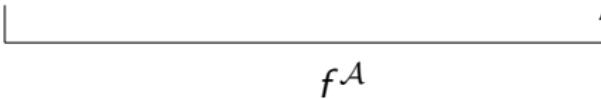
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Indeed, $f^{\mathcal{A}}(\text{free}(n)) = h(\langle n, i \rangle)$ and i is greater than the position of $\text{free}(n)$.

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We construct a new punctual copy \mathcal{B} in the following way:

$$\tilde{c}(0) = 0, \quad \tilde{c}(2^{i_k} + 2^{i_{k-1}} + \dots + 2^{i_0}) = 2^{c(i_k)} + 2^{c(i_{k-1})} + \dots + 2^{c(i_0)},$$

so that the individuals of \mathcal{B} have the form $\tilde{c}(n)$ and

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Idea: we can simulate binary addition by looking ahead with the successor $S^{\mathcal{A}}$ in the original model \mathcal{A} .

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Proposition

Let $+$ be primitive recursive in the model \mathcal{A} . Let $p(x) = 2^x$.
Then $p^{\mathcal{B}}$ is primitive recursive if and only if $p^{\mathcal{A}}$ is primitive recursive.

Method of islands and archipelago

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At stage e we have the following picture:

mainland $0 \rightarrow a_1 \rightarrow a_2 \rightarrow a_3 \rightarrow \dots \rightarrow a_k$

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In this stage we wait for $p_e(w)$ to give value v for s steps.

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Method of islands and archipelago (2)

While waiting we must keep extending the mainland with new elements and also the archipelago with new islands.

mainland $0 \rightarrow a_1 \rightarrow a_2 \rightarrow a_3 \rightarrow \dots \rightarrow a_k \rightarrow a_{k+1} \rightarrow a_{k+2}$

archipelago $w^x \quad b_1^{\tau_1} \quad b_2^{\tau_2} \quad \dots \quad b_m^{f(\tau_1)} \quad b_{m+1}^{g(\tau_2)} \quad b_{m+2}^{f(\tau_2)}$

When we obtain the result $p_e(w) = v$ we must connect the mainland with the archipelago.

How to connect?

We must choose carefully a position q for w , because this choice will fix the positions of all islands.

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But since the terms over \mathcal{F} should be closed under substitution, we omit the base- x case from the definition.

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After choosing q , every island b_i obtains the corresponding position, which is its label τ_i evaluated at q .

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We also add other auxiliary elements, so that the mainland and the islands become one successor chain:

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And finally, if it happens that $S^A(v) = w$ we must insert a new element a between them, because we want to ensure that $\text{pred}^A(w) \neq v$ (so that $p_e \neq \text{pred}^A$).

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We pick a new element w' , which is the start of a new archipelago and we proceed to stage $e + 1$.

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Thanks for your attention!

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