

Learning Rational Probability Distributions with Bisequential Variational Autoencoders

Georgi Shopov

Probabilistic Models

L'essentiel est invisible pour les yeux

$$\mathbb{P}_S$$

The essential is invisible to the eyes

French sentence S

English translation



$$\mathbb{P}_U$$

There's no smoke without fire

Utterance U

Transcription

A delicate and precise watercolor illustration ...

$$\mathbb{P}_D$$



Text description D

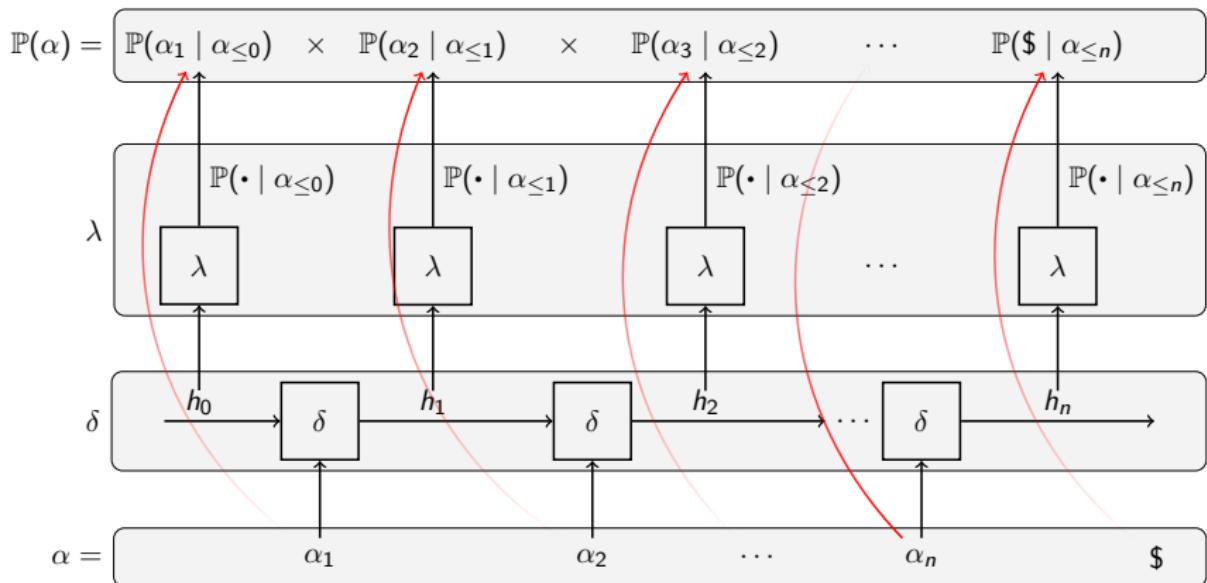
Image

Recurrent Probabilistic Models

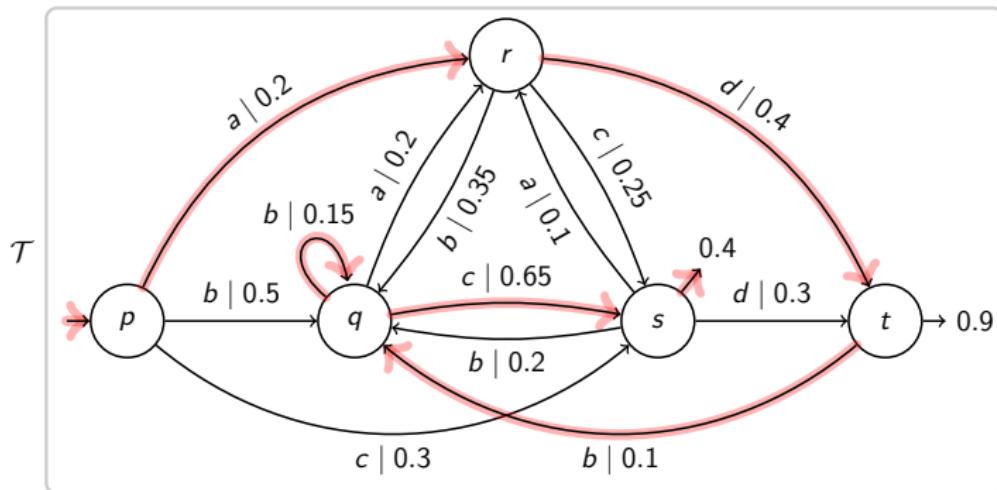
$$h_0 \in \mathbb{R}^d$$

$$\delta: \mathbb{R}^d \times \Sigma \rightarrow \mathbb{R}^d$$

$$\lambda: \mathbb{R}^d \rightarrow \Delta^{|\Sigma|}$$



Sequential Transducers



$$\llbracket T \rrbracket(adbabc) = 0.2 \times 0.4 \times 0.1 \times 0.15 \times 0.65 \times 0.4$$

Sequential transducers \rightarrow sequential functions

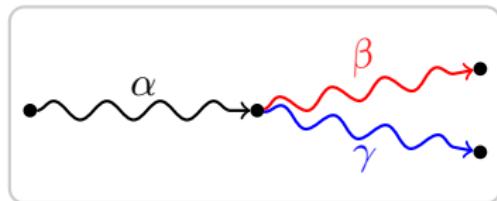
Arbitrary transducers \rightarrow rational relations

Expressive Power of Sequential Transducers

Definition

For $\alpha, \beta, \gamma \in \Sigma^*$ such that $\beta \wedge \gamma = \epsilon$, the **prefix distance** between $\alpha\beta$ and $\alpha\gamma$ is defined as

$$d_p(\alpha\beta, \alpha\gamma) = |\beta| + |\gamma|.$$



Theorem (Mohri [3])

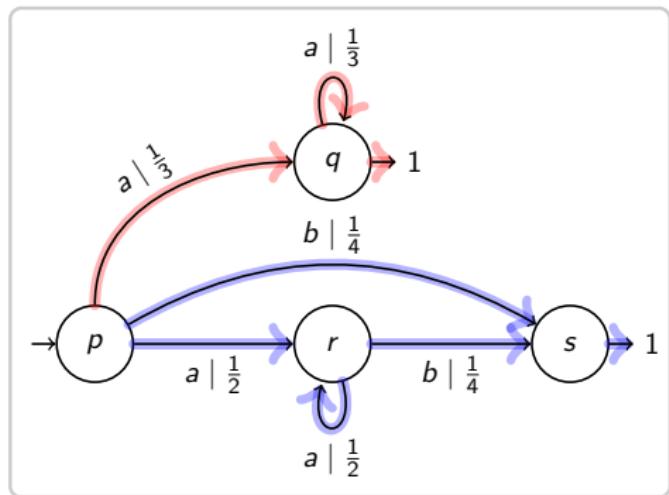
A rational probability distribution \mathbb{P} over Σ^* is sequential if and only if

$$\left\{ \frac{\mathbb{P}(\alpha)}{\mathbb{P}(\beta)} \mid \alpha, \beta \in \text{Supp}(\mathbb{P}) \text{ & } d_p(\alpha, \beta) \leq n \right\}$$

is finite for all $n \in \mathbb{N}$.

Limitations of Sequential Transducers

$$\mathbb{P}(\alpha) = \begin{cases} \left(\frac{1}{3}\right)^{n+1} & \text{if } \alpha = a^{n+1} \\ \left(\frac{1}{2}\right)^{n+2} & \text{if } \alpha = a^n b \\ 0 & \text{otherwise} \end{cases}$$

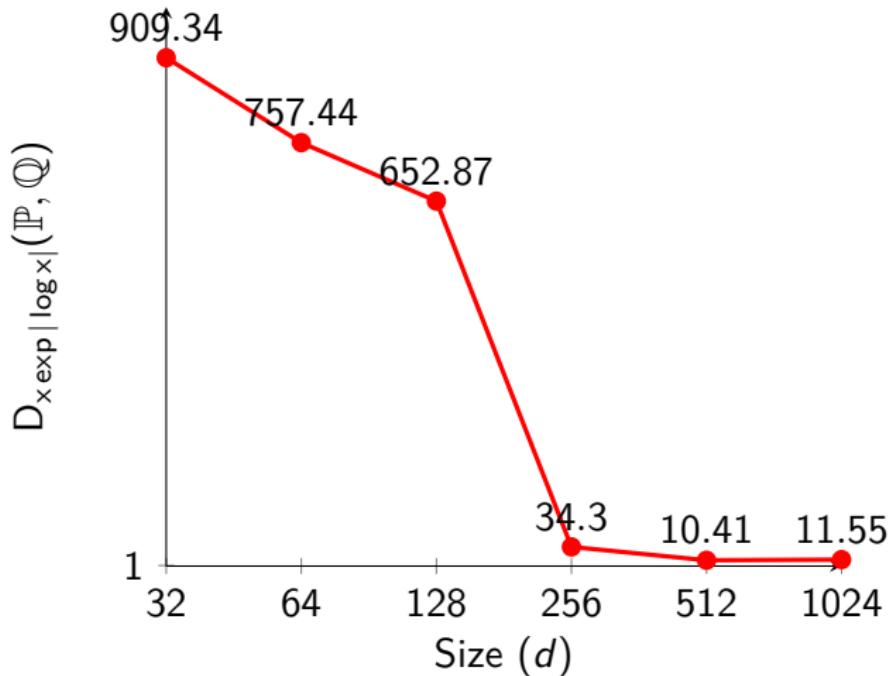


\mathbb{P} is rational but not sequential because for every $n \in \mathbb{N}$,

$$d_p(a^{n+1}, a^n b) = 2 \quad \text{and} \quad \frac{\mathbb{P}(a^n b)}{\mathbb{P}(a^{n+1})} = \frac{\left(\frac{1}{2}\right)^{n+2}}{\left(\frac{1}{3}\right)^{n+1}} = \frac{1}{2} \left(\frac{3}{2}\right)^{n+1}.$$

Limitations of Recurrent Probabilistic Models

$$D_{\exp|\log x|}(\mathbb{P}, \mathbb{Q}) = \mathbb{E}_{\alpha \sim \mathbb{P}} \left[\exp \left| \log \frac{\mathbb{P}(\alpha)}{\mathbb{Q}(\alpha)} \right| \right] = \mathbb{E}_{\alpha \sim \mathbb{P}} \left[\frac{\max\{\mathbb{P}(\alpha), \mathbb{Q}(\alpha)\}}{\min\{\mathbb{P}(\alpha), \mathbb{Q}(\alpha)\}} \right]$$

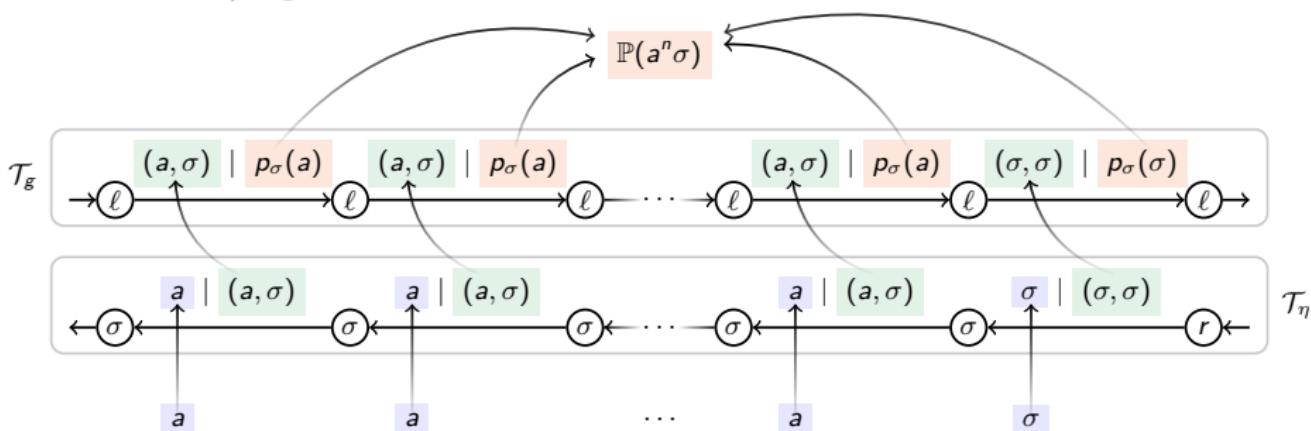


Bisequential Decompositions

Definition

A **bisequential decomposition** of a probability distribution \mathbb{P} over Σ^* is a tuple (Γ, η, g) , where

- Γ is a **latent alphabet**;
- $\eta: \Sigma^* \rightarrow (\Sigma \times \Gamma)^*$ is a co-sequential function s.t. $\eta \circ \pi_{\Sigma^*} = \text{id}_{\Sigma^*}$;
- $g: (\Sigma \times \Gamma)^* \rightarrow [0, 1]$ is a sequential probability distribution;
- $\mathbb{P} = \eta \circ g$.



Expressive Power of Bisquential Decompositions

Theorem (Elgot and Mezei [1])

A probability distribution is rational if and only if it admits a bisquential decomposition.

Theorem (Shopov and Gerdjikov [5])

A probability distribution \mathbb{P} over Σ^ is rational if and only if there exists a finite partition $\{L_i\}_{i=1}^n$ of Σ^* such that, for $1 \leq i \leq n$, L_i is regular and $\{\mathbb{P}(\cdot \alpha \mid L_i \alpha)\}_{\alpha \in \Sigma^*}$ is finite. In this case,*

$$\alpha \sim_i \beta \iff \mathbb{P}(\cdot \alpha \mid L_i \alpha) = \mathbb{P}(\cdot \beta \mid L_i \beta)$$

is a left congruence and the latent alphabet can be chosen to be

$$\Gamma = \left\{ \left(\mathbb{P}(\cdot \alpha \mid L_i \alpha) \right)_{i=1}^n \mid \alpha \in \Sigma^* \right\}.$$

Latent Variable Models

- A latent variable model specifies a joint distribution

$$p_{\theta}(x, z)$$

over observed variables (data) $x \in \mathcal{X}$ and latent variables $z \in \mathcal{Z}$.

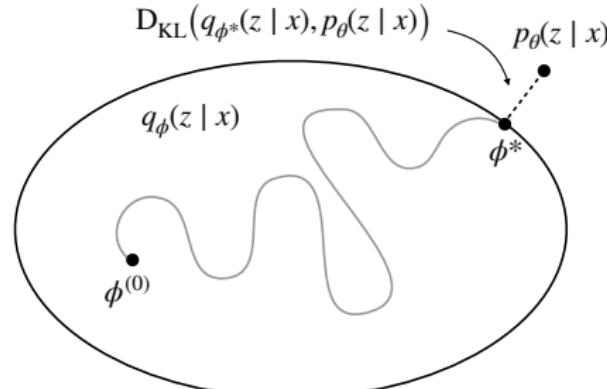
- The latent variables explain the ***hidden structure used to generate the data.***
- However, the marginal distribution over the data

$$p_{\theta}(x) = \int_{z \in \mathcal{Z}} p_{\theta}(x, z) dz$$

is often intractable.

Evidence Lower Bound

$$\begin{aligned}\log p_\theta(x) &= \mathbb{E}_{q_\phi(z|x)} [\log p_\theta(x)] \\ &= \mathbb{E}_{q_\phi(z|x)} \left[\log \frac{p_\theta(x, z) q_\phi(z|x)}{p_\theta(z|x) q_\phi(z|x)} \right] \\ &= \mathbb{E}_{q_\phi(z|x)} \left[\log \frac{p_\theta(x, z)}{q_\phi(z|x)} \right] + \mathbb{E}_{q_\phi(z|x)} \left[\log \frac{q_\phi(z|x)}{p_\theta(z|x)} \right] \\ &= \underbrace{\mathbb{E}_{q_\phi(z|x)} \left[\log \frac{p_\theta(x, z)}{q_\phi(z|x)} \right]}_{\text{Evidence Lower Bound}} + \underbrace{\mathbb{D}_{\text{KL}}(q_\phi(z|x), p_\theta(z|x))}_{\geq 0}\end{aligned}$$



Variational Autoencoders

Definition (Kingma and Welling [2], Rezende et al. [4])

A **variational autoencoder** is a pair (ϕ, θ) of parameters defining

- a variational encoder $q_\phi(z | x)$, and
- a generative model $p_\theta(x, z) = p_\theta(x | z)p_\theta(z)$

that are optimised by maximising the evidence lower bound

$$\begin{aligned}\mathbb{E}_{q_\phi(z|x)} \left[\log \frac{p_\theta(x, z)}{q_\phi(z|x)} \right] &= \mathbb{E}_{q_\phi(z|x)} \left[\log \frac{p_\theta(x|z)p_\theta(z)}{q_\phi(z|x)} \right] \\ &= \mathbb{E}_{q_\phi(z|x)} [\log p_\theta(x|z)] + \mathbb{E}_{q_\phi(z|x)} \left[\log \frac{p_\theta(z)}{q_\phi(z|x)} \right] \\ &= \underbrace{\mathbb{E}_{q_\phi(z|x)} [\log p_\theta(x|z)]}_{\text{Reconstruction term}} - \underbrace{\mathbb{D}_{\text{KL}}(q_\phi(z|x), p_\theta(z))}_{\text{Regularisation term}}.\end{aligned}$$

Bisequential Variational Autoencoders

Definition

A variational autoencoder is called **bisequential** if

- $q_\phi(z | x) = \prod_{i=1}^n q_\phi(z_i | x_{\geq i})$, and
- $p_\theta(x, z) = \prod_{i=1}^n p_\theta(z_i | x_{<i}, z_{<i}) p_\theta(x_i | x_{<i}, z_{\leq i})$.

Theorem

The evidence lower bound of a bisequential variational autoencoder (ϕ, θ) can be expressed as

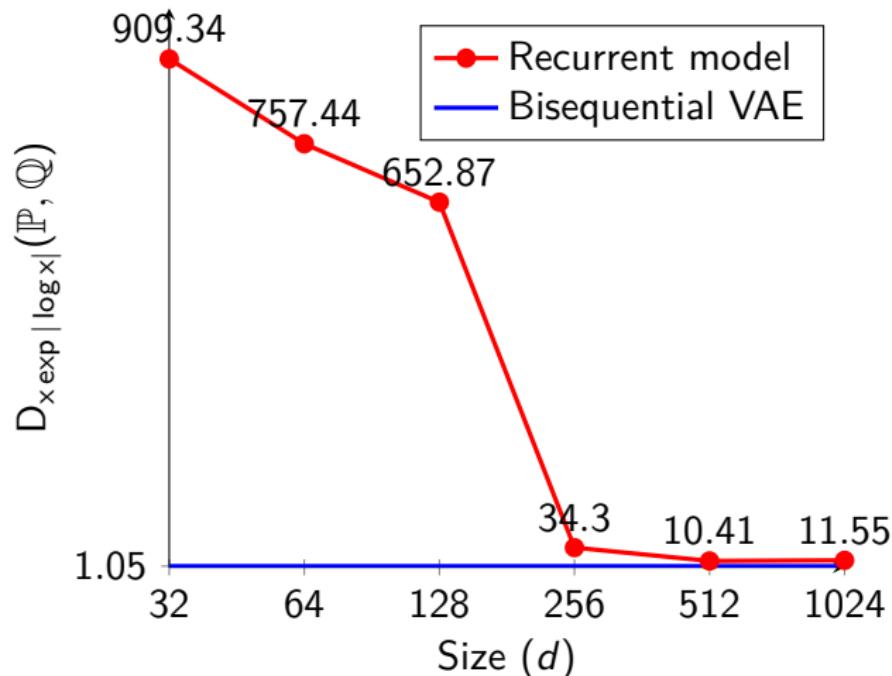
$$\mathbb{E}_{q_\phi(z|x)} \left[\sum_{i=1}^n \underbrace{\log p_\theta(x_i | x_{<i}, z_{\leq i})}_{\text{Reconstruction term}} - \underbrace{D_{\text{KL}}(q_\phi(z_i | x_{\geq i}) || p_\theta(z_i | x_{<i}, z_{<i}))}_{\text{Regularisation term}} \right].$$

Furthermore, when approximating a rational probability distribution, this lower bound will be tight.

Expressive Power of Bisequential VAEs

Theorem

Every rational probability distribution can be represented by a bisequential variational autoencoder.



Conclusion

We demonstrated how automata theory can be used to:

- ***Characterise the expressive power of recurrent models.***
- ***Identify key limitations*** – in particular, their inability to represent ***non-sequential rational probability distributions.***
- ***Motivate more expressive architectures***, such as ***Bisequential VAEs***, which overcome these limitations and can model ***the full class of rational probability distributions.***

References I

- [1] C. C. Elgot and J. E. Mezei. On relations defined by generalized finite automata. *IBM Journal of Research and Development*, 9(1):47–68, 1965.
- [2] Diederik P. Kingma and Max Welling. Auto-encoding variational bayes. In *The Second International Conference on Learning Representations*, 2014.
- [3] Mehryar Mohri. Finite-state transducers in language and speech processing. *Computational Linguistics*, 23(2):269–311, 1997.
- [4] Danilo Jimenez Rezende, Shakir Mohamed, and Daan Wierstra. Stochastic backpropagation and approximate inference in deep generative models. In Eric P. Xing and Tony Jebara, editors, *Proceedings of the Thirty-First International Conference on Machine Learning*, volume 32 of *Proceedings of Machine Learning Research*, pages 1278–1286. PMLR, June 2014.

References II

[5] Georgi Shopov and Stefan Gerdjikov. Consistent bidirectional language modelling: Expressive power and representational conciseness. In Yaser Al-Onaizan, Mohit Bansal, and Yun-Nung Chen, editors, *Proceedings of the 2024 Conference on Empirical Methods in Natural Language Processing*, pages 5724–5768. Association for Computational Linguistics, November 2024.